Code-Level Verification for Autonomy (TC2)

Robust and Resilient Autonomy for Advanced Air Mobility

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Credits: NASA / Lillian Gipson







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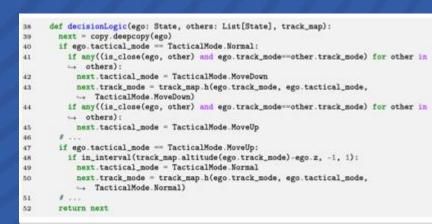


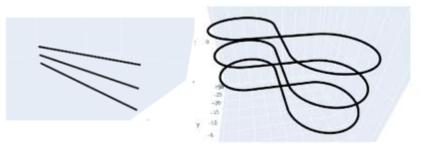
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Outline

Introduction to Verse verification tool https://github.com/AutoVerse-ai/Verse-library

- Quadrotor example
- Demo

Background theory: Data-driven verification

Reachability and sensitivity

Applications

- L1 Adaptive control [Lin Song et al., ICCPS WIP '23]
- DNN control [Puthumanaillam, Ornik, et al.]
- RL air-traffic management [Peng Wei, GWU, ongoing]
- Parallel Verse [Zhu, et al.]

Verification problem and the Verse tool

Input

System + Requirements Multi-agent hybrid scenarios + Invariance / safety

Output

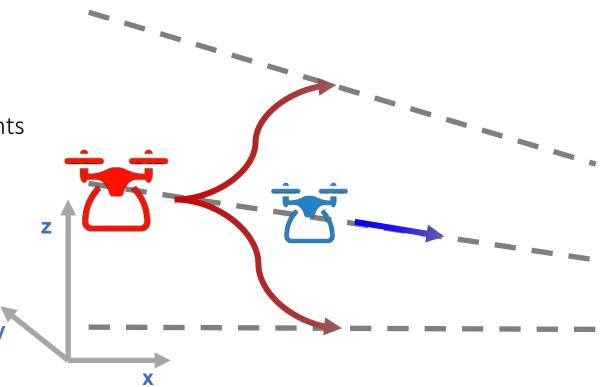
Counter-example or Proof that all behaviors of the System meet Requirements

Basic Verse Approach

Probabilistic sensitivity analysis + Deterministic reachability analysis

Advanced Approaches

Parallelized verification Incremental verification Verification with neural network controllers





Creating scenarios in Verse --- quick and easy

Drones flying along three parallel vertically separated straight tracks

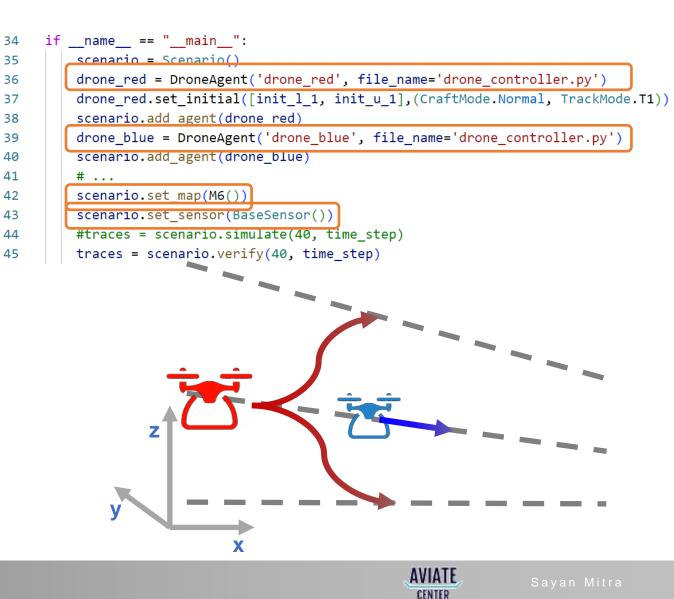
Scenario in Verse defined by a set of agents, a map, and a sensor

Maps define sequences of waypoints or motion primitives

Sensors specify information available to one agent about other agents

Agent: *Decision logic* + *Dynamics*

- 6D model + Bang-bang controller
- 18D model + L1 controller

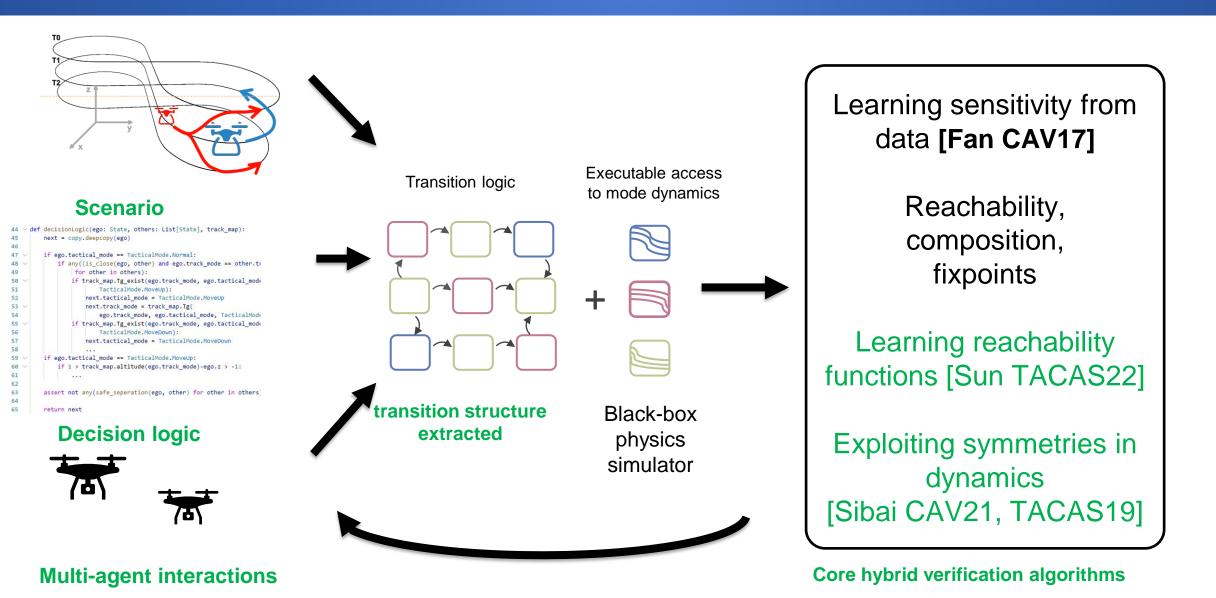


Writing Decision Logic

<pre>44 ~ def decisionLogic(ego: State, others: List[State], track_map): 45</pre>	Informal logic: "Nondeterministically switch to track above or below if too close to any other drone in the same track"
if thatk_map.rg_exist(ego.thatk_mode, ego.tatticat_mode, (Decision Logic (DL) modifies agent's mode written in Python
55 V (if thack man Tg exist(ego thack mode ego tactical mode)	If exists multiple possible transitions, both will be explored
<pre>58 59 \sigma if ego.tactical_mode == TacticalMode.MoveUp: 60 \sigma if 1 > track_map.altitude(ego.track_mode)-ego.z > -1: 61</pre>	Using <i>any</i> and <i>all</i> to quantify over other agents in the scenario
<pre>62 63 63 64 64</pre> 62 63 64	DL supports user defined functions
	Safety specified using python <i>asserts</i> (no need to learn new logics)



Verse under the hood



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Verse scenario to hybrid system

Consider a scenario SC with k agents,

Agent i has:

- $_{\circ}$ $\,$ Continuous state space X_{i}
- $\circ \quad \text{Mode space } D_i$
- $_{\odot}~$ Guard G_i and reset R_i extracted from Decision Logic. For a pair of modes $d,d' \in D_i$
 - $G_i(d, d') \subseteq \prod_j X_j$ defines transition condition
 - $R_i(d, d'): \prod_j X_j \rightarrow X_i$ defines change of continuous state after transition
- $_{\circ}~$ Flow function $F_i {:}~ X_i \times D \times \mathbb{R}^{\geq 0} \rightarrow X_i$



Verse scenario to hybrid system cont.

Hybrid system $H(SC) = \langle X, X^0, D, d^0, G, R, TL \rangle$ from composition of agents Continuous state space $X = \prod_k X_i$

◦ An element $x \in X$ is called a state; X^0 : set of initial states

Mode space $D = \prod_k D_i$

• An element $d \in D$ is called a mode. d^0 : an initial mode

For a pair of modes $d, d' \in D$

- ∘ guard G(d, d') ⊆ X, x ∈ G(d, d') iff x ∈ $G_i(d_i, d'_i)$ and $d_j = d'_j$ for $j \neq i$
- reset $R(d, d'): X \to X, R(d, d')(X) = \prod_k R_i(d, d')(X)$

TL is a set of pairs $\langle \xi, d \rangle$

- Trajectory $\xi: [0, T] \rightarrow X$ describes evolution of continuous states in mode $d \in D$
- Given $d \in D$, $x \in X$, $\forall t, \xi(t) = \prod_j F_j(x_j, d_j)(t)$
- ξ . *fstate*, ξ . *lstate*, ξ . *ltime* the first state $\xi(0)$, the last state of the trajectory $\xi(T)$, and ξ . *ltime* = T

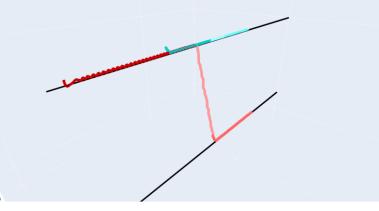
Execution and reachable states for H(SC)

 δ -execution of H(SC) is a sequence of m labeled trajectories $\alpha \coloneqq \langle \xi^0, d^0 \rangle, ..., \langle \xi^{m-1}, d^{m-1} \rangle$

- \circ ξ⁰.fstate ∈ X⁰
- $\circ \ \xi^{i-1}. \ \text{lstate} \in G(d^{i-1}, d^i) \ \text{and} \ \xi^i. \ \text{fstate} = R\big(d^{i-1}, d^i\big)(\xi^{i-1}. \ \text{lstate}), \ i \in \{1, m-1\}$
- \circ ξⁱ.ltime = δ if i ≠ m − 1 and ξ^{m−1}.ltime ≤ δ, i ∈ {1, m − 1}

Reach (x⁰, d⁰, T_{max}): reachable states from x⁰ ∈ X, d⁰ ∈ D along an execution α defined as $U_{i \in [0,m)} U_{t \in [0,\delta)} \xi^{i}(t), \text{ with } m = \left[\frac{T_{max}}{\delta}\right]$ $\circ \text{ Reach}_{H}(T_{max}) = U_{x^{0} \in X^{0}} \text{ Reach } (x^{0}, d^{0}, T_{max})$

To prove safety, we can check $Reach_H(T_{max}) \cap Unsafe = \emptyset$





For a pair of modes d, d', define discrete and continuous post operators

- postCont(X, d) = X' iff X' = $\bigcup_{x \in X} \prod_K F_i(x_i, d, \delta)$
- postDisc(X, d, d') = X' iff $\forall x \in X, x \in G(d, d')$ and $X' = \bigcup_{x \in X} R(d, d')(x)$

Verse constructs reachability tree Tree = $\langle V, E \rangle$ up to depth m

- Each vertex (S, d) ∈ V is a pair of a set of continuous states and a mode
- $_{\circ}$ Root $\langle X^{0},d^{0}\rangle$
- There is an edge from (S, d) to (S', d'), iff S' = postCont(postDisc(S, d, d'), d')

Reachability tree constructed by Verse is an over-approximation of $Reach_H(T_{max})$



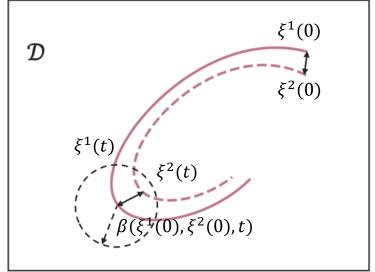
PostCont: Using Discrepancy Function

A discrepancy function $\beta : \mathbb{R}^n \times \mathbb{R}^n \times \mathbb{R}^{\geq 0}$ is a uniformly continuous function such that for any pair of trajectories $\langle \xi^1, d \rangle, \langle \xi^2, d \rangle \in TL$, and any $t \in \xi^1$. dom $\cap \xi^2$. dom

- $\circ |\xi^{1}(t) \xi^{2}(t)| \leq \beta(\xi^{1}. \text{fstate}, \xi^{2}. \text{fstate}, t)$
- \circ β(.,.,t) → 0 as ξ^1 .fstate → ξ^2 .fstate

Compute postCont for input set of states $X^0 = Ball(x^0, r)$ in mode d

- $\,\circ\,\,$ Obtain trajectory ξ^0 starting from x^0 labeled by mode d
- $\circ~$ Obtain discrepancy function β
- An over-approximation of reachable set can be obtained by $U_t \xi^0(t) \bigoplus \beta(x_0, x_0 + r, t)$





Learn discrepancy from data

Use a template for exponential discrepancy $\beta(x_1, x_2, t)$

 $||x_1(t) - x_2(t)|| \le \beta(x_1, x_2, t) = ||x_1(0) - x_2(0)||e^{at+b}$

Taking log:

 $\forall t, \ln \frac{\|x_1(t) - x_2(t)\|}{\|x_1(0) - x_2(0)\|} \le at + b$

Find *a* and *b* by **learning a linear separator**

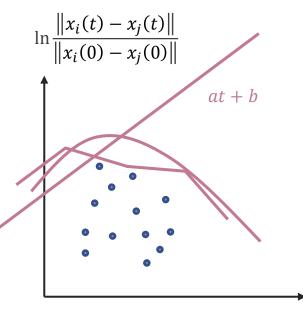
Theorem [CAV17]: Given the training set, the global exponential discrepancy function that gives the tightest reach set over-approximation can be found by solving a Linear programming (LP) problem

Proposition [CAV17]: $\forall \epsilon, \delta > 0$, if sampling number $n \ge \frac{1}{\epsilon} \ln \frac{1}{\delta}$, then with probability $1 - \delta$, the algorithm finds (a, b) such that $err_{\mathcal{D}}(a, b) < \epsilon$



AVIATE

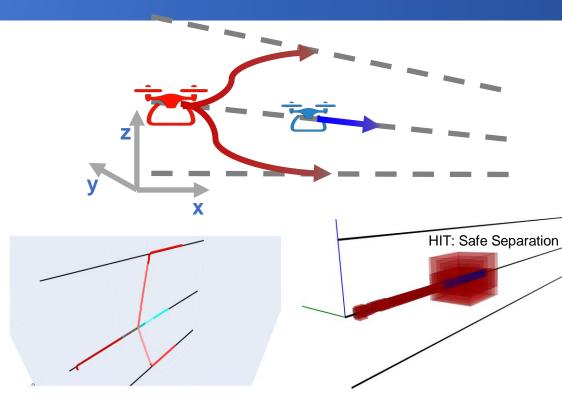
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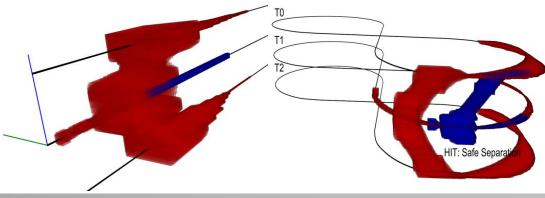


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Result highlights

- Computed trajectory of 2-drone scenario
- Reachability analysis find potential safety violation
- By Modifying parameters in the decision logic, we can mitigate the safety violation
- Easily modify scenario to test more interesting behaviors of the agents
 - Further show with live demo







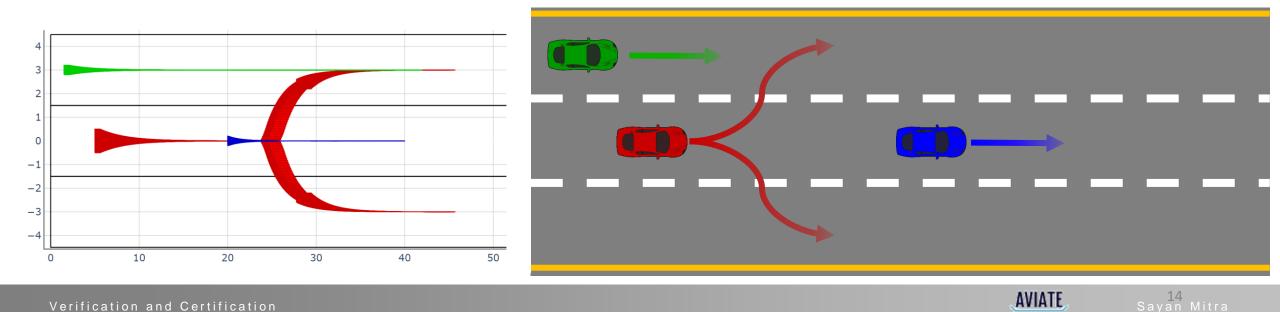
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Live Demo1: Lane Switch Scenario

Three cars 0

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- Red and green cars running at 1 m/s and blue cars running at 0.5m/s
- Red car can perform lane switch when there's another car 5m in front 0
- Three vehicles all have uncertain initial condition
- Safety condition: The red vehicle should be 1m away from all other vehicles.



CENTER



Live Demo2: Easy Modification Detect Safety Violation

 Exact same setting as base example but with a different Map

60	<pre>car = CarAgent('car1', file_name=input_code_name)</pre>	
61	<pre>scenario.add_agent(car)</pre>	
62	<pre>car = NPCAgent('car2')</pre>	
63	<pre>scenario.add_agent(car)</pre>	
64	<pre>car = NPCAgent('car3')</pre>	
65	<pre>scenario.add_agent(car)</pre>	
66	<pre>tmp_map = M1()</pre>	
67	<pre>scenario.set_map(tmp_map)</pre>	
60 61 62 63 64 65 66	<pre>car = CarAgent('car1', file_name=input_code_name) scenario.add_agent(car) car = NPCAgent('car2') scenario.add_agent(car) car = NPCAgent('car3') scenario.add_agent(car) tmp_map = M3() scenario.set_map(tmp_map)</pre>	

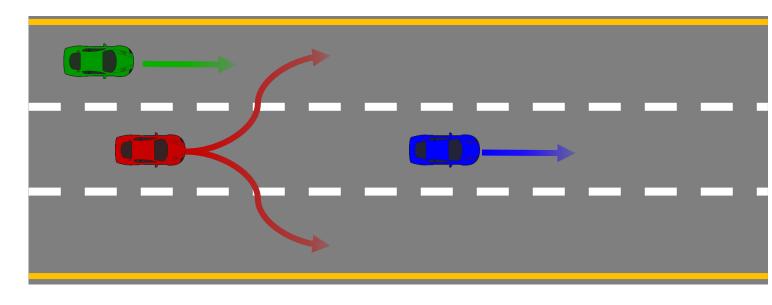


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Live Demo3: Handle Uncertainty in Perception

 Exact same setting as base example but with different sensor model with noise

60	<pre>car = CarAgent('car1', file name=input code name)</pre>
61	scenario.add agent(car)
62	<pre>car = NPCAgent('car2')</pre>
63	scenario.add agent(car)
64	<pre>car = NPCAgent('car3')</pre>
65	<pre>scenario.add agent(car)</pre>
66	tmp map = M1()
67	scenario.set_map(tmp_map)
37	<pre>config = ScenarioConfig(init_seg_length=5)</pre>
38	scenario = Scenario()
39	<pre>car = CarAgent('car1', file_name=input_code_name)</pre>
40	<pre>scenario.add_agent(car)</pre>
41	<pre>car = NPCAgent('car2')</pre>
42	<pre>scenario.add_agent(car)</pre>
43	<pre>car = NPCAgent('car3')</pre>
44	<pre>scenario.add_agent(car)</pre>
45	$tmp_map = M1()$
46	scenario.set map(tmp map)
47	<pre>scenario.set_sensor(NoisyVehicleSensor(\</pre>
48	(0.5, 0.5), (0.0, 0.0)))



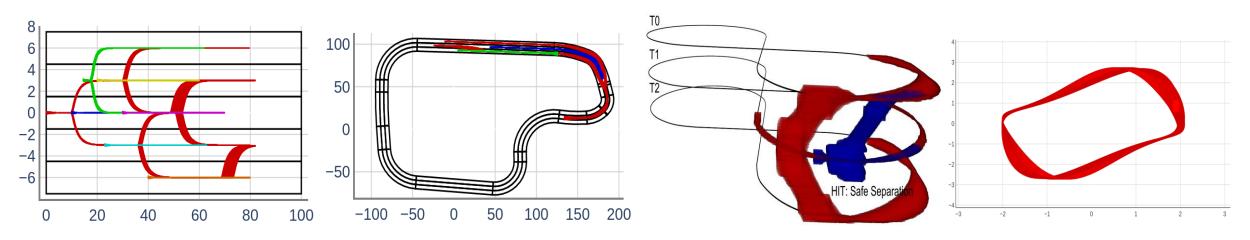




More Scenarios Verified by Verse

Table 1: Runtime for verifying examples in Section 5. Columns are: number of agents (# \mathcal{A}), agent type (\mathcal{A}), map used (Map), reachability engine used (postCont), sensor type (Noisy \mathcal{S}), number of mode transitions #TR, and the total run time (Rt). N/A for not available.

$\#\mathcal{A}$	$ \mathcal{A} $	Map	postCont	Noisy \mathcal{S}	#Tr	Rt (s)	$\#\mathcal{A}$	\mathcal{A}	Map	postCont	Noisy \mathcal{S}	#Tr	Rt (s)
2	D	$\mathcal{M}6$	DryVR	No	8	55.9	2	D	$\mathcal{M}5$	DryVR	No	5	18.7
2	D	$\mathcal{M}5$	NeuReach	No	5	1071.2	3	D	$\mathcal{M}5$	DryVR	No	$\overline{7}$	39.6
7	C	$\mathcal{M}2$	DryVR	No	37	322.7	3	С	$\mathcal{M}1$	DryVR	No	5	23.4
3	C	$\mathcal{M}3$	DryVR	No	4	34.7	3	С	$\mathcal{M}4$	DryVR	No	$\overline{7}$	118.3
3	C	$\mathcal{M}1$	DryVR	Yes	5	29.4	2	С	$\mathcal{M}1$	DryVR	No	5	21.6
2	C	$\mathcal{M}1$	NeuReach	No	5	914.9	1	V	N/A	DryVR	N/A	1	0.33
1	S	N/A	DryVR	N/A	3	2.3	1	G	N/A	DryVR	N/A	3	67.14

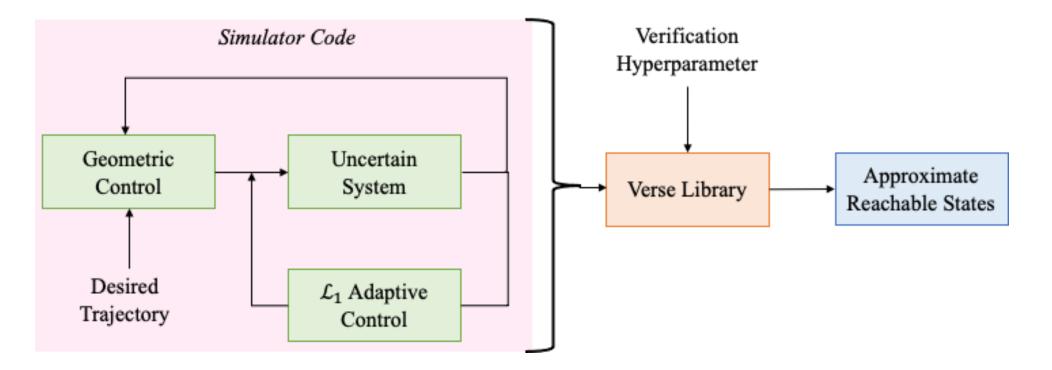




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 L₁ Adaptive Control (L1AC) verification architecture using the Verse Library [Song et al. ICCPS-WIP 23]



L. Song, Y. Li, S. Cheng, P. Zhao, S. Mitra, N. Hovakimyan, Verification of \$\mathcal{L}_1\$ Adaptive Control using Verse Library: A Case Study of Quadrotors, *In Proceedings of* 13th IEEE International Conference on Cyber Physical Systems (ICCPS) Demo/Poster/Work-in-Progress, San Antonio, TX, 2023.

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L1AC verification objectives

Formally verify the following two properties of L1AC:

• Transient performance guarantees;

- Scenario: an 18-dimensional drone model subject to rapidly changing uncertainty
- Expected outcome: L1AC's capability for fast adaptation

• Guaranteed delay margins.

- Scenario: an 18-dimensional drone model subject to time delay in the control input
- Expected outcomes:
 - L1AC preserves delay margin bounded away from zero;
 - Graceful performance degradation provided by L1AC.

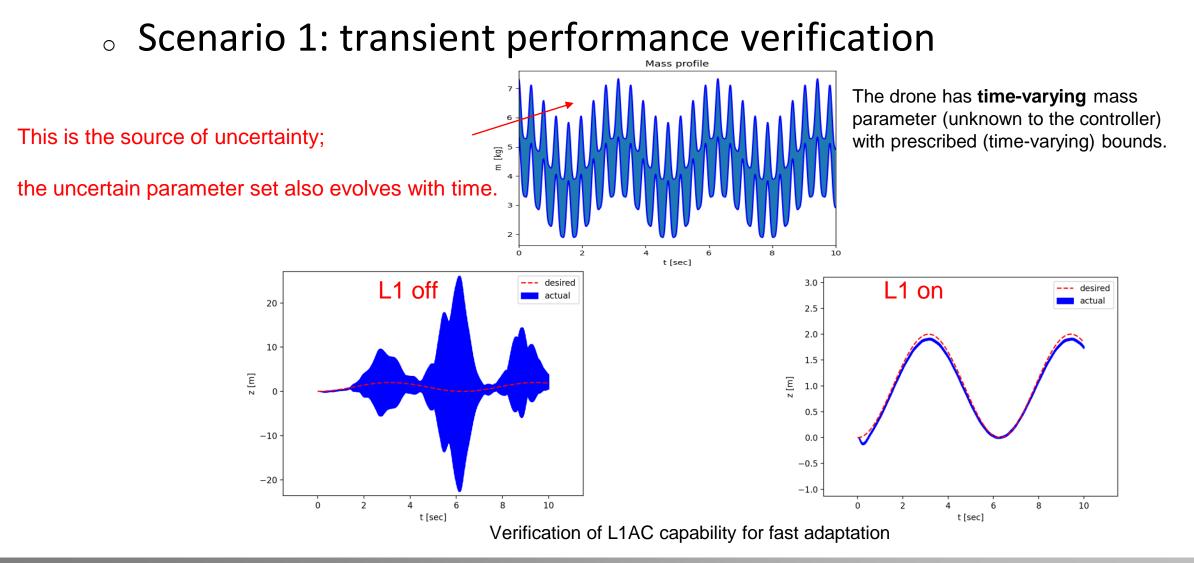


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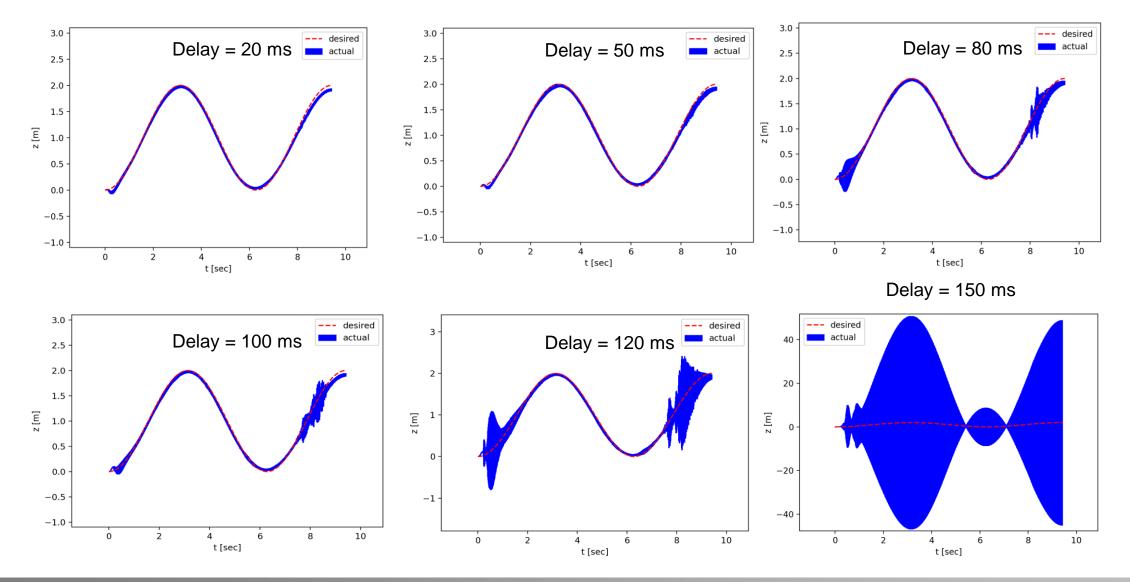
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Scenario 2: delay margin verification

Control input of the drone system is subject to **<u>time delay</u>**.

Only consider the delay margin achieved by L1AC, and we implement the verification procedure under **a range of** time delay amount.





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• Future directions:

- Verification of L1AC plus learning-enabled component;
 - An example: Contraction L1AC + Gaussian Processes [Gahlawat et al. L4DC 2021]
- Verification of L1AC on systems involving switch, either on model or controller;
 - An example: learn-to-fly [Snyder et al. JGCD 2022], vehicle subject to driving environment changes [Mao et al. ACM TCPS 2023]
 - Tool/Method: deploy the mode switch feature of the Verse Library
- Verification of the controller-parameter tuning process.
 - An example: Difftune⁺ [Cheng et al. L4DC 2023]
 - Tool/Method: Postdisc + Postcont -- 'one-step' reachability analysis feature of Verse



Other applications

 Application to DNN-based control [Puthumanaillam, Ornik, et al.]

 Application to RL-based air-traffic management [Peng Wei, GWU, ongoing]

• Parallel Verse [Zhu, et al.]



Summary

- Verse is designed to make hybrid system verification accessible
 - Python DL, nondeterministic agents, scenarios, sensors, asserts, OpenDrive maps
- Under the hood Verse uses tree-based reachability, sensitivity analysis for postCont
 - Can handle uncertainty in initial states, transitions, parameters
 - Plug-in Post computations DryVR, NeuReach, Monotonicity, ...
- $_{\circ}$ In the future
 - Incremental verification, parallelization
 - verifying DRL controllers
- We welcome your feedback!

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Thank You Very Much!



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